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# Random thoughts on random walks:

networks, centrality, and tracking the spread of disease

### networks

- fancy word for graphs: nodes = vertices,
   connections = edges
- typically directed (Twitter) or undirected (Facebook)
- connectedness/giant component, irreducibility



### adjacency matrix

- 1 if edge from *i* to *j* (0 otherwise)
- symmetric for undirected graphs
- makes counting easier
  - degree of node *i*:  $d_i = \sum_j a_{ij}$
  - # of walks of length t starting from node i:

$$M_i(t) = \sum_{i_1, i_2, \dots, i_t} a_{ii_1} a_{i_1 i_2} \cdots a_{i_{t-1} i_t}$$

# of 3-cycles in entire network (triple counting):

$$C^{(3)} = \sum_{i} (A^3)_{ii} = Tr(A^3)$$







# what does it mean to be **central** within a network? context-dependent!

create your own...

<u>http://vax.herokuapp.com/game</u>







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# centrality

• • •

- {in-, out-, undirected-} degree centrality "How many {followers, followees, friends} do I have?"
- eigenvector centrality "How influential am I?"
- Katz centrality/PageRank "If my friends are important, doesn't that make me important?"
- closeness centrality "How close am I to everyone else, on average?"
- betweenness centrality "How many chains of connections include me?"

# degree centrality

$$c_i^{deg} = d_i / \sum_i d_i$$

- normalization doesn't matter
- defined using only local information (up to scaling)

### eigenvector centrality

defined recursively/iteratively: requires global knowledge of network



 $A\mathbf{v} = \kappa \mathbf{v}$ 





similar to eigenvector centrality, but using normalized adjacency matrix 

defined using global information (as before)

- add a damping factor to ensure everyone ends up with at least some centrality
  - i.e., interpolate original graph with complete graph

$$a_{ij} = \frac{a_{ij}}{d_i}$$

 $P\pi = \pi$ 

Page et. al (1999)



# (a disservice to) Markov chains

transition probabilities of going from *i* to *j* 

s $\begin{array}{c}
s_1 \\
\vdots \\
s_r \\
p_{r1}
\end{array}$ 

- stationary distribution satisfies  $\int \pi_i p_i$
- rate of convergence given by second largest eigenvalue of transition matrix
- example: simple random walk on finite set of states

$$\begin{pmatrix} 1 & \cdots & s_r \\ 1 & \cdots & p_{1r} \\ & & & \vdots \\ 1 & \cdots & p_{rr} \end{pmatrix}$$

$$_{ij}=\pi_j$$
 (i.e.,  $\pi P=\pi$ )

# PageRank + random walks

- interpretation: proportion of time random walker spends at each node, assuming at each step neighbor selected at random
  - i.e., stationary distribution of Markov process on graph with uniform transition probabilities
  - technical details: irreducibility, aperiodicity, ergodicity
  - damping factor ensures these properties hold
- note for later: only local information used at each step in walk...

## degree centrality + random walks

- for undirected networks, PageRank (w/o damping) is the simple random walk
- has degree centrality as its stationary distribution if network is connected
  - check this yourself! show that

 $\sum_{i} p_{ij} \left( \frac{a_j}{\sum_i d_i} \right) = \frac{a_i}{\sum_i d_i}$ 

### degree centrality + random walks

- moments of degree distribution:  $\langle d^m \rangle$
- if we sample a node uniformly at random: P(X = i) = 1/N =
- according to their degree centralities:

$$P(X=i) = d_i / \sum_i d_i \qquad \Longrightarrow \qquad$$

intuition: In the limit, every edge is traversed same proportion of times. So, we are sampling a node at the end of an edge chosen uniformly at random!

$$= \frac{1}{N} \sum_{i} d_{i}^{m} = \sum_{d} d^{m} P(d)$$

$$\Rightarrow E(deg(X)) = \langle d \rangle$$

instead, after performing random walk "long enough", we are sampling nodes

$$E(deg(X)) = \frac{\langle d^2 \rangle}{\langle d \rangle} = \langle d \rangle + \frac{Var(d)}{\langle d \rangle}$$



### sampling using random walks

- constructive: If we want to sample according to degree centrality, take any node and select neighbor uniformly at random... rinse, repeat...
- principle behind Markov Chain Monte Carlo (MCMC)
  - provides way to generate samples from distribution that can't be sampled directly
  - need to specify transition probabilities such that stationary distribution of Markov chain is the desired sampling distribution

# real-world implications

- Christakis & Fowler (2010)
  - flu outbreak among Harvard student population from Sep.-Dec. '09
  - 2009 H1N1 pandemic: ~60M infected in US, ~300K hospitalizations
  - two groups: random vs. friends of random (FoR)
  - epidemic curve in FoR tracked progression in random group ~14d in advance
  - deal with it: on average, your friends have more friends than you

Strogatz, New York Times (2012)





- long history of using (random) walks to infer properties of the social graph
- Milgram's small-world experiment (1967)
  - Watts-Strogatz model: shortcuts in highly clustered networks
  - Kleinberg model: distribution of shortcut lengths and efficient routing



### Erdos-Renyi (tree-like)



### Erdos-Renyi (tree-like) small diameter, no clustering



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### Watts-Strogatz (lattice + shortcuts)



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# small diameter, high clustering Watts-Strogatz (lattice + shortcuts)





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### Kleinberg (distribution of shortcuts)





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Kleinberg (distribution of shortcuts) greedy routing finds shortcuts!





### Erdos-Renyi (tree-like) small diameter, no clustering

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Watts-Strogatz (1998); Kleinberg (2000)

### disease surveillance

- what do public health officials and CDC care about?
  - situational awareness
  - early detection of epidemic onset
  - peak timing and intensity
- practical reasons
  - vaccine supply and distribution
  - allowing hospitals to run at high capacity and prepare for large influx of patients
- difficulties both mathematical/statistical and practical



# let's do a little math

SIR model on a graph (recall Andy's talk!)

$$\frac{ds_i}{dt} = -\beta \sum_j a_{ij} s_i x_j$$
$$\frac{dx_i}{dt} = \beta \sum_j a_{ij} s_i x_j - \gamma x_i$$
$$\frac{dr_i}{dt} = \gamma x_i$$



Newman (book); message passing: Karrer-Newman (2013)





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bi



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$$\frac{dr_i}{dt} = \gamma x_i$$

### for small times, P(i is S) is approximately 1

 $v_i$ 





Newman (book); message passing: Karrer-Newman (2013)





# linearized SIR on graph

 $\frac{d\mathbf{x}}{dt} = \beta A\mathbf{x} - \gamma \mathbf{x}$ 

# linearized SIR on graph

 $\mathbf{x}(t) = \exp(t(A - \gamma I))\mathbf{x}(0) \approx e^{(\beta\kappa - \gamma)t}\mathbf{v} \quad \mathbf{v}$ 

# $\frac{d\mathbf{x}}{dt} = \beta A\mathbf{x} - \gamma \mathbf{x}$

### dominant e-value of A

# linearized SIR on graph

# $\frac{d\mathbf{x}}{dt} = \beta A\mathbf{x} - \gamma \mathbf{x}$

### dominant e-value of A

 $\mathbf{x}(t) = \exp(t(A - \gamma I))\mathbf{x}(0) \approx e^{(\beta\kappa - \gamma)t}\mathbf{v} \quad \mathbf{v}$ 

epidemic takes off if  $R_0 = \frac{\beta}{\gamma} > \frac{1}{\kappa}$ 

dominant e-value of A

 $\mathbf{x}(t) = \exp(t(A - \gamma I))\mathbf{x}(0) \approx e^{(\beta\kappa - \gamma)t}\mathbf{v} \quad \mathbf{v}$ 

$$\mathbf{x}(t) = \exp(t(A - \gamma I))$$

$$p = \frac{\mathbf{x}(\tau_S) \cdot \mathbf{1}_S}{M} = \frac{e^{(\beta \kappa - \gamma)\tau_S} \mathbf{v} \cdot \mathbf{1}_S}{M}$$

expected prevalence hits p in surveillance set

dominant e-value of A

 $\mathbf{\mathbf{x}}(0) \approx e^{(\beta\kappa - \gamma)t} \mathbf{v}$ 

 $p = \frac{\mathbf{x}(\tau) \cdot \mathbf{1}}{N} = \frac{e^{(\beta \kappa - \gamma)\tau} \mathbf{v} \cdot \mathbf{1}}{N}$ 

expected prevalence hits p in entire network

$$\mathbf{x}(t) = \exp(t(A - \gamma I))$$

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expected prevalence hits p in surveillance set

$$\Delta \tau := \tau_S - \tau = \frac{1}{\beta \kappa - \gamma} \ln\left(\frac{c}{c_S}\right)$$

dominant e-value of A

 $\mathbf{\mathbf{x}}(0) \approx e^{(\beta\kappa - \gamma)t} \mathbf{v}$ 

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 $p = \frac{\mathbf{x}(\tau) \cdot \mathbf{1}}{N} = \frac{e^{(\beta \kappa - \gamma)\tau} \mathbf{v} \cdot \mathbf{1}}{N}$ 

expected prevalence hits p in entire network

average e-vector centrality in surveillance set

average e-vector centrality in network

- problem: not possible\* without knowing entire network to begin with!
- never really know the underlying social graph, can only infer global properties from local subnetworks
- is there a random walk whose stationary distribution is eigenvector centrality?



### correlation of centrality measures



- could use degree centrality (locally recoverable) as a proxy for eigenvector centrality
- can we do better: is there a random walk whose stationary distribution is eigenvector centrality?



$$p_{ij} = \frac{a_{ij}v_j}{\kappa v_i}$$

transition probabilities



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transition probabilities

 $\overline{\psi_i} = v_i^2$ 

### stationary distribution!



$$p_{ij} = \frac{a_{ij}v_j}{\kappa v_i}$$

transition probabilities



$$\psi_i = v_i^2$$

### stationary distribution!

$$\sum_{j} p_{ij} \psi_j = \frac{1}{\kappa} v_i \sum_{j} a_{ij} v_j = v_i^2 = \psi_i$$



$$p_{ij} = \frac{a_{ij}v_j}{\kappa v_i}$$

transition probabilities



$$\psi_i = v_i^2$$

### stationary distribution!

$$\sum_{j} p_{ij} \psi_j = \frac{1}{\kappa} v_i \sum_{j} a_{ij} v_j = v_i^2 = \psi_i$$

but transition probabilities given in terms of e-vector centralities!



### maximal entropy random walk

- why maximal entropy?
- uniform distribution on a set has maximal entropy among all distributions on that set
- transition probabilities of MERW put equal probability on all **paths** of length t starting from a given node, as t goes to  $\infty$





### entropy rate

 $i_1,...,i_t$ 

# probability of particular path $S(t) = -\sum p(i, i_1, \dots, i_t) \ln p(i, i_1, \dots, i_t)$

### Entropy of set of paths of length t starting at i

### entropy rate

$$S(t) = -\sum_{i_1,\dots,i_t} p(i,$$

### Entropy of set of paths of length t starting at i

h =

entropy rate

## probability of particular path $(i_1,\ldots,i_t)\ln p(i,i_1,\ldots,i_t)$

$$\lim_{t \to \infty} \frac{\ln S(t)}{t}$$

### entropy rate

minimal entropy: all probability is put on one path starting from *i*:

$$S(t) = -\sum_{i_1,...,i_t} p(i, i_1, ..., i_t) \ln p(i, i_1, ..., i_t) = 0$$

maximal entropy: uniform probability on all paths of length t starting from i:

$$S(t) = -\sum_{i_1,...,i_t} \frac{1}{M_i(t)} \ln\left(\frac{1}{M_i(t)}\right) = \ln M_i(t)$$

maximal entropy

$$M_i(t) = \sum_{i_1, i_2, \dots, i_t} a_{ii_1} a_{i_1 i_2} \cdots a_{i_{t-1} i_t}$$

# 0 for all paths except one!

### 

uniform probability on all paths

$$p(i, i_1, \dots, i_t) = \frac{a_{ii_1} a_{i_1 i_2} \dots a_{i_{t-1} i_t}}{\sum_{i_1, i_2, \dots, i_t} a_{ii_1} a_{i_1 i_2} \dots a_{i_{t-1} i_t}}$$

uniform probability on all paths

$$\pi(i_1|i) = \lim_{t \to \infty} \frac{a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}{\sum_{i_1} a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}$$

corresponding transition probabilities!

 $\pi(i_1|i) = \lim_{t \to \infty} \frac{a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}{\sum_{i_1} a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}$ 



$$\pi(i_1|i) = \lim_{t \to \infty} \frac{a_{ii_1} \sum_{i_1} \sum_{i_1} a_{ii_1}}{\sum_{i_1} a_{ii_1}}$$

$$\pi^0(i_1|i) = \frac{a_{ii_1}}{\sum_{i_1} a_{ii_1}}$$

**Oth-order** approximation "How many friends do I have?"



 $\frac{\sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}{\sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}$ 

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$$\pi^0(i_1|i) = \frac{a_{ii_1}}{\sum_{i_1} a_{ii_1}}$$

**Oth-order** approximation "How many friends do I have?"



$$\pi^{1}(i_{1}|i) = \frac{a_{ii_{1}}k(i_{1})}{\sum_{i_{1}}a_{ii_{1}}k(i_{1})}$$

**1st-order** approximation "How many friends do my friends have?"

$$\pi(i_1|i) = \lim_{t \to \infty} \frac{a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}{\sum_{i_1} a_{ii_1} \sum_{i_2} a_{i_1 i_2} \dots \sum_{i_t} a_{i_{t-1} i_t}}$$

$$\pi^0(i_1|i) = \frac{a_{ii_1}}{\sum_{i_1} a_{ii_1}}$$

### **Oth-order** approximation "How many friends do I have?"

$$\pi^{2}(i_{1}|i) = \frac{a_{ii_{1}}\sum_{i_{2}}a_{i_{1}i_{2}}k(i_{2})}{\sum_{i_{1}}a_{ii_{1}}\sum_{i_{2}}a_{i_{1}i_{2}}k(i_{2})}$$
$$= \frac{a_{ii_{1}}k(i_{1})k_{nn}(i_{1})}{\sum_{i_{1}}a_{ii_{1}}k(i_{1})k_{nn}(i_{1})},$$



$$\pi^{1}(i_{1}|i) = \frac{a_{ii_{1}}k(i_{1})}{\sum_{i_{1}}a_{ii_{1}}k(i_{1})}$$

### **1st-order** approximation "How many friends do my friends have?"



**2nd-order** approximation "What is the average degree of my friends' friends?"

# approximated MERW



# performance on Montreal network





- can use random walks to sample nodes with desired properties
- how to incorporate into a realistic sampling methodology?
- network models are fun to work with!

