

Drawbacks with triangulated categories! as basic but not as good as Hilbert spaces.

DG & A_{infty} category ... probably equivalent richer structures, also simpler, more naturally in derived category context.

k commutative rings

Complex = complex of k -modules

usual picture of complex... alternatively $\mathcal{C} = \bigoplus \mathcal{C}^i$, $d: \mathcal{C} \rightarrow \mathcal{C}$ $\deg \mathcal{C} = i$, $d^2 = 0$:

i.e. \mathcal{C} is a DG k -module : better POV psychologically!

complexes form a tensor category (symmetric monoidal category)

$$\otimes: \{k\text{-DG-mod}\} \times \{k\text{-DG-mod}\} \longrightarrow \{k\text{-DG-mod}\}$$

simplest to write in DG language:

$\mathcal{C} \otimes \mathcal{C}'$ usual \otimes on underlying k -modules with sum of gradings, differential using (gradings) Leibniz formula

... think uniformly in language of \otimes categories rather than in "complex"-picture.

$(A \otimes B) \otimes C \xrightarrow{\sim} A \otimes (B \otimes C)$, $A \otimes B \xrightarrow{\sim} B \otimes A$ functorial isomorphisms satisfying identities.

Associativity for complexes obvious, commutativity is supercommutativity, $a \otimes b \xrightarrow{k} (-)^{p+q} b \otimes a$ $p = \deg a$, $q = \deg b$

$k = \text{unit of } \otimes$.

DG - category = category enriched over k -DG-modules.

A category has objects $Ob \mathcal{A}$,

& $\text{Hom}(X, Y) \in k\text{-DG-mod}$ morphism

composition $\text{Hom}(X, Y) \otimes \text{Hom}(Y, Z) \rightarrow \text{Hom}(X, Z)$ associative.

Identity: $X \in Ob \mathcal{A} \Rightarrow$ morphism $k \rightarrow \text{Hom}(X, X)$

map $k \xrightarrow{\sim} \text{id}_X$ element $z \in \mathcal{C}$ with $\deg z = 0$ & $dz = 0$

so $\text{id}_X \in \text{Hom}(X, X)$ $\deg 0$ closed.

Can consider ~~of~~ plain category: forget grading & differentials
 $\rightarrow k$ -linear category

Another categorical way to pass from enriched to usual category: can replace \mathcal{C} by set $\text{Hom}(k, \mathcal{C}) =$ cocycles of degree zero ...

Examples a. $k\text{-DG-mod}$ is a DG-category:

$\text{Hom}(\mathcal{C}, \mathcal{C}')$ has natural DG-mod: grading comes from natural \mathbb{G}_m action on Hom ...

replace whole Hom by sum of homogeneous submodules (n-graded mod) $\text{Hom}(\mathcal{C}, \mathcal{C}')^n = \prod_{k=0}^n \text{Hom}(\mathcal{C}^k, \mathcal{C}'^{k+n})$

Morphism of complexes is a 0-cocycle in this dg hom $\text{Hom}(\mathcal{C}, \mathcal{C}')$ i.e. $\text{Hom}_{\text{complex}}(\mathcal{C}, \mathcal{C}') = \text{Hom}(k, \text{Hom}(\mathcal{C}, \mathcal{C}'))$

So DG-categories are categories with dg k -mod structure on hom sets so that compositions are k -linear.

Where does commutativity of \otimes come in? \rightsquigarrow

can speak of dual category of DG category \mathcal{A}°

Example 1 DG category with one object

\longleftrightarrow DG-algebra (initial - implicit herefore)

2. A usual k -linear category = DG category with Hom s are dg \mathbb{O} , $d=0$.

3. Complexes of R -modules, R an algebra

3'. Complexes of R -modules, R a k -linear category:

R -module = k -linear functor $\mathcal{R} \rightarrow k\text{-mod}$

- psychologically important that \mathcal{R} are exactly k -containing...

e.g. left A -mod & right A -mod \Rightarrow group/bundles

- some construction for R -modules

e.g. $X_1, \dots, X_n \in \mathcal{R}$ for any objects set each object of \mathcal{R} is isomorphic to a direct sum (finite) of the X_i
 $\rightarrow R\text{-module is isomorphic to } \text{End}_R \bigoplus X_i$ formal sum.

Even if R arbitrary can form formal direct sum $\bigoplus_{X,Y} \text{Hom}(X, Y)$

- get algebra but without unit, only very idempotents

4. DG-modules over a DG-algebra $R = \text{Alg}_k$

(Q: special case of 3' for $R = \mathbb{C}[[t]]$)

Another POV: a quasifactor is a bimodule over A_1, A_2 with certain properties.

Another POV: $A\text{-factors}$ of the DG category ...

- i.e. replace A by a certain (co)unit (counit) morphism
- at least over a field!

Triangulated Categories

Example $K(C)$: homotopy category of complexes in a k -linear category

A triangulated category is a graded category with additional structure: A "candidate triangle" is

$$\deg f = \deg g = 0, \deg h = 1, gf = 0 = hs = fh$$

$$\begin{array}{ccc} & f \rightarrow Y & g \\ X & \xrightarrow{h} & Z \end{array}$$

Additional structure: some candidate Δ 's are "distinguished" + axioms...

Don't require shift as structure: only in this graded setting ask for representability of suspension:

Action 0 $\forall X \in \mathcal{T} \quad \forall n \in \mathbb{Z} \quad \exists Y \xrightarrow{f} X \quad \deg f = n$
s.t. f^{-1} exists.

\Rightarrow defined up to unique isomorphism, call $Y = \Sigma^n X$
 \Rightarrow weak action of \mathbb{Z} on \mathcal{T} .

$K(C)$: distinguished Δ is given by $X \xrightarrow{f} Y \xrightarrow{\text{Core } f} \text{Co}(f \rightarrow X)$

$$(\text{Core } f)^n = Y^n \otimes X^{n+1}$$

$$dx[n] = -dx$$

$$d((n)) = \begin{pmatrix} d & f \\ 0 & -dx \end{pmatrix}$$

Another POV (edge): $\dots \rightarrow Y^{n-1} \xrightarrow{f} Y^n \xrightarrow{\text{Core } f} Y^{n+1} \rightarrow \dots$ d_{Co}

$$X^{n-1} \rightarrow Y^n \xrightarrow{\text{Core } f} X^{n+1} \rightarrow X^{n+2}$$

$$d_{\text{Co}}$$

$$d_{\text{Co}}$$

- forms a double complex (chain), & core is total complex ... really naive double complex!

$$df = fd \quad \text{no signs}$$

don't set $(d' + d'')^2 = 0$, so must introduce signs

Examples a. $K(C)$ or $K(R\text{-mod})$, R a k -algebra

b. $H_0(R\text{-DG mod})$, R a DG algebra or

c. $H_0(R\text{-DG mod})$

T a triangulated category $\rightarrow T'$ a full subcategory.
 a candidate triangle in T is distinguished if it is so in T'
 ... ask these to give Δ structure on $T' \rightarrow$ triangulated subcategory
 $\hookrightarrow T'$ closed wrt suspensions, desuspensions & cones
 in weak sense

Advantage of DG categories: Yoneda $A \xrightarrow{\sim} A^{\circ}$ -DG-mod
 Triangulated are have $T \hookrightarrow_{\text{tautolog}} T \rightarrow$ graded modulos
 and in fact $T \hookrightarrow$ cohomological functors, i.e. graded functors
 $T^{\circ} \rightarrow$ Graded bimodules. Higher cohomological functors
 don't have natural triangulated structure.

DG categories II

A DG-category \rightarrow homotopy category $Ho^*(A)$

$A \rightarrow A^{\circ}$, DG functor is a quasi-equivalence & induces
 equivalence $Ho^*(A) \rightarrow Ho^*(A^{\circ})$.

would like to localize "world" of DG categories wrt quasi-equivalences.

Triangulated category: graded category + additional structure, class of distinguished triangles.

Yoneda: $A \xrightarrow{\text{full}} A^{\circ}$ -DG-mod, $Ho^*(A) \hookrightarrow Ho^*(A^{\circ}$ -DG-mod)

Def. A candidate triangle in $Ho^*(A)$ has triangulated structure if it
 is distinguished in $Ho^*(A^{\circ}$ -DG-mod)
 A is said to be pre-triangulated if $Ho^*(A)$ is
 triangulated after equipped with this structure, i.e.
 if $Ho^*(A)$ is a triangulated subcategory of $Ho^*(A^{\circ}$ -DG-mod),
 — i.e. closed under cones & desuspensions.

Problem: the def is not self-contained: what to get two nodes
 of distinguished triangle, opposite to each other, not clear
 if they agree. Do DG functors preserve these
 distinguished triangles? \rightsquigarrow more covariant to reformulate correctly

4' DG-motives over a DG-category R $\& \text{DG-nat}$
 $(\text{DG-functor } R \rightarrow k\text{-DG-mot})$

For any DG-category R has Yoneda embedding $R \hookrightarrow R^0 \text{-dg-mot}$
(right R -nat) $X \mapsto h_X$, contravariant
 $h_X(Y) = \text{Hom}(Y, X)$

"embedding": fully faithful $(h_X(Y) \xrightarrow{\sim} [\text{Hom}(F(X), F(Y))] \text{ isom})$

A -DG category, $\text{Ho}(A) = \underline{\text{homotopy category}}$.

Or $\text{Ho}(A) = \text{dg } A$

$X, Y \in \text{dg } A$ $\text{Ho}(A)$ -morphisms $X \rightarrow Y$ are $H^0 \text{Hom}(X, Y)$.

Other notation $\text{Ho}(A) = H^0(A)$.

$\text{Ho}^0(A)$ = graded category, the graded homotopy category, with
morphisms elements of $\bigoplus H^n \text{Hom}(X, Y) \equiv \text{Ext}^n(X, Y)$
-- can have property (not strict) of being
triangulated -- if not can add two objects to make triangulated.

Def A DG-functor $F: A_1 \rightarrow A_2$ is a quasi-equivalence
if $H^0(F): \text{Ho}^0(A_1) \rightarrow \text{Ho}^0(A_2)$ is an equivalence, i.e.
1. $\text{Hom}(X, Y) \xrightarrow{\sim} \text{Hom}(F(X), F(Y))$ is a quism.
2. exactness property of $\text{Ho}(F): \text{Ho}(A_1) \rightarrow \text{Ho}(A_2)$
(may graded version): $\forall Z \in A_2 \exists X \in A_1$,
 $f: F(X) \rightarrow Z$ $\deg f = 0$ $df = 0$, f has
a homotopy inverse.

A_1, A_2 DG-categories are quasi-equivalent if
 $\begin{array}{ccc} A_1 & \xrightarrow{\sim} & A_2 \\ \downarrow & \cong & \downarrow \\ A_1 & \xrightarrow{\sim} & A_2 \end{array}$ diagram of quasi-equivalences,

- Not enough to define comonadic or "weak"
of dg categories ... as equivalence to quasi-equivalences.
-- what do dg functors form?
"Funzy" notion of quasi-functor $A_1 \rightarrow A_2$:

$\begin{array}{ccc} A_1 & \xrightarrow{\sim} & A_2 \\ \downarrow & \text{dg-functor} & \downarrow \\ \text{quasi-equivalence} & \xrightarrow{\sim} & \text{dg-functor} \end{array}$ These form (allowing Kortz) a DG-category
in Ass setting

Need for $A^\circ\text{-DG-mod}$: carries cones of morphisms from \mathcal{A}° , $\text{Cone}(f: X \rightarrow Y) = \text{Cone}(h_X \xrightarrow{f} h_Y)$
 on representable functors
 — don't need full $A^\circ\text{-DG-mod}$ — just cones (& iterated ones)
 & desuspensions, so don't need big $A^\circ\text{-DG-mod}$.

\mathcal{C} pre-additive category (have some of morphisms but not all direct sums) — to formally add direct sums can go several routes:

$$\mathcal{C} \hookrightarrow \mathcal{C}^\circ\text{-mod}$$

Or concretely just look at

category with objects $\bigoplus \mathcal{C}_i$ (both have finite \oplus)
 formal direct sums, morphisms just matrices (rather than "linear transformations" in other defns).

"Def of
"pretriangulated
full"

Step 1 Replace $A^\circ\text{-DG-mod}$ by $\tilde{\mathcal{A}} \subset A^\circ\text{-DG-mod}$

full subcategory: smallest full DG-subcategory s.t. $\tilde{\mathcal{A}} \supseteq \mathcal{A}$ and

• $M \in \tilde{\mathcal{A}}$, $M' \cong M$ isomorphism $\Rightarrow M' \in \tilde{\mathcal{A}}$

• $M \in \tilde{\mathcal{A}}$, $n \in \mathbb{Z} \Rightarrow M[n] \in \tilde{\mathcal{A}}$.

• $M \Rightarrow$ stable wrt semisplit extensions: $0 \rightarrow M_1 \rightarrow M_2 \rightarrow M_3 \rightarrow 0$

is semisplit extension (of DG-modules) if split in sense of graded modules eg cone $0 \rightarrow N \rightarrow \text{Cone}(M \rightarrow N) \rightarrow M[1] \rightarrow 0$

- f denotes failure of splitting to be closed -

its differential is f .

semisplitting

More concretely $M \in \tilde{\mathcal{A}} \iff \exists \text{ finite filtration } 0 = M_0 \subset \dots \subset M_n = M$

finite filtration s.t. $M_i/M_{i-1} \cong h_{a_i}[\tau_i]$ $a_i \in A, \tau_i \in \mathbb{Z}$

- representable up to shift. ($\text{char } A^\circ\text{-DG-mod} \leftrightarrow \text{char } A$)

-- the extensions $0 \rightarrow M_{i+1} \rightarrow M_i \rightarrow M_i/M_{i-1} \rightarrow 0$ are automatically semisplit because representable functors are projective objects

$h_a \in A^\circ\text{-graded mod}$.

eg suppose $A \supseteq A$ is an algebra (ie only one object)

- as A ! object \Rightarrow module $h_A = A$ as A -module, which is free \Rightarrow projective.

Step 2 Fix the splittings $0 \rightarrow M_{i+1} \rightarrow M_i \rightarrow M_i/M_{i-1} \rightarrow 0$

So $M = \bigoplus h_{a_i}[\tau_i]$ as graded module.

As a DG-module $M = (\bigoplus_{i=0}^r A_{[r]} \otimes_{\mathbb{Z}} \mathbb{F}_0^{\oplus d} + \omega)$

d = standard differential on direct sum of dg modules

$\omega = (\omega_{ij})$ $\omega_{ij} \in \text{Hom}(A_j, A_i)[r_i - r_j]$
of degree 1 (wrt grading)

1. ω is strictly upper triangular $\omega_{ij} = 0$ if $i > j$
2. $D_\omega^2 = 0$ i.e. $D_\omega \omega + \omega^2 = 0$ flat connection.
(Maurer-Cartan eqn).

Def A^{prefr} the pretriangulated hull of A : "DG category of twisted complex"
 $\mathcal{O}_b A^{\text{prefr}}$ are formal expressions $(\bigoplus_{i=0}^r A_{[r]}, \omega)$
 $\omega = (\omega_{ij})$ as above.

$\text{Hom}(\bigoplus A_{[r]}, \bigoplus A'_{[r']})$ as graded module
(ie ignore d , since ω) is just space of matrices

$f = (f_{ij})$ $f_{ij} \in \text{Hom}(A_j, A'_i)[r_i - r_j]$.

composition automatic
 $D f = D_\omega f - f D_\omega (-1)^{r_j}$ super commutator:
symbol D has degree 1. = change + $c_0 f - (-1)^{\deg f} f \omega$.

'twisted complexes' --- Toledo-Tay.

This definition is clearly self-dual, & clear that
DG functors extend automatically to pretriangulated hulls.

$f: X \rightarrow Y$ in A $\Rightarrow \text{Core}(f) \in A^{\text{prefr}}$ canonical object
 $\text{Core}(f) = (Y \otimes X_{[r]}, (\circ f))$.
 $X \in A$, $X_{[r]} \in A^{\text{prefr}}$

Def $X \xleftarrow{\quad} \xrightarrow{\quad} Y$ in $\text{Ho}(A)$ is distinguished if it is
isomorphic to $X \xleftarrow{f} \xrightarrow{\quad} \text{Core}(f)$

A is pretriangulated if $\forall f: X \rightarrow Y$ in A , $\text{Core}(f)$
is homotopy equivalent to an object of A , & see for $X_{[r]}$ no \mathbb{Z} .

Exercise 1. Notion of distinguished Δ is self dual

2. $F: A \rightarrow B$ DG-functor $\Rightarrow \text{Ho} F: \text{Ho} A \rightarrow \text{Ho} B$

Sends distinguished triangles to distinguished triangles. If Ho^{f} fully faithful \rightarrow converse also holds.

3. A^{pretr} is a pretriangulated DG-category
4. $\text{Ho}^{\text{f}}(A^{\text{pretr}})$ is generated as Δ -category by $\text{Ho}^{\text{f}}(A)$
5. A is pretriangulated iff $A \rightarrow A^{\text{pretr}}$ is a quasi-equivalence.
6. $A^{\text{pretr}} \rightarrow (A^{\text{pretr}})^{\text{pretr}}$ is an equivalence of DG-categories.
7. If $F: A \rightarrow B$ is a quasi-equivalence then it induces a quasi-equivalence $F^{\text{pretr}}: A^{\text{pretr}} \rightarrow B^{\text{pretr}}$

Bondal-Kapranov notation $A^{\text{tr}} := \text{Ho}^{\text{f}}(A^{\text{pretr}})$.

A DG-category \rightarrow triangulated category.

Do derived categories, e.g. $D(A\text{-mod})$ or $D(\text{abelian dg})$ come from a DG-category? 2 abelian categories are complicated objects ... take another view

A a DG category $\Rightarrow D(A) :=$ derived category
of A -modules $= \text{Ho}^{\text{f}}(A^{\text{D-mod}})$ (exact DG-mod)
- quotient of triangulated category by subcategory.

Importance of this example (B. Keller): have

Yoneda $A \xleftarrow{\text{inj}} A^{\text{D-mod}}$ extends \Rightarrow fully faithful
 $A^{\text{tr}} := \text{Ho}^{\text{f}}(A^{\text{pretr}}) \hookrightarrow D(A)$
fully faithful --- so any of the A^{tr} 's embed into
such a derived category of modules.

Exercise*: A quasi-equivalence $A \rightarrow B$ induces an equivalence $D(A) \rightarrow D(B)$: one restriction factor is an adjunction $\text{Res} \dashv \text{LInd}$ of derived functor of induction

$D(A) \xrightleftharpoons[\text{Res}]{\text{LInd}} D(B)$ always adjoint, & in this case colimitors are isomorphisms.

Q: $D(A) = \text{Ho}^{\text{f}}(?)$? DG-category containing A .
Answer (known to topologists through question next):
- Spaltenstein, Arramir - Halperin, Hirsch

Will define $\xrightarrow{A} \rightarrow A^0\text{-dg-mod}$: category of
semi-free DG-modules (in cell complexes ...)
so that $H_0(\xrightarrow{A}) \rightarrow D(A)$ is an equivalence.

Well known (in bounded setting) : R a usual ring,
{Derived category of bounded above} = {homotopy category of bounded above
complexes of R -modules} = {complexes of projective R -modules}

try to eliminate boundedness:

Example $R = \mathbb{Z}/4\mathbb{Z}$ (assumed that R has ∞ homological dimension)
 $\dots \rightarrow R \xrightarrow{2} R \xrightarrow{2} R \xrightarrow{2} R \rightarrow \dots$
acyclic complex but not homotopic to zero:
apply $\otimes_{\mathbb{Z}} \mathbb{Z}/2\mathbb{Z}$, preserves null-homotopy, but get
 $\dots \mathbb{Z}/2 \xrightarrow{0} \mathbb{Z}/2 \xrightarrow{0} \mathbb{Z}/2 \rightarrow \dots$
with nonzero cohomology!

So to eliminate boundedness need projectivity assumption on
whole complex not its terms..

Exercise If A has finite homological dimension (\exists finite proj. resoln.)
then can remove boundedness in above ~~at the~~ equivalence.
(in particular acyclic complexes of projectives will be null-homotopic)

Def A DG- A^0 -module P is free if it is DG-isomorphic
to $\bigoplus_i \text{Im } [e_i]$ (possibly infinite sum)
 \implies projective
(in case of just a DG-algebra! module generated by e_i is of
degree $-r$, with $de_i = 0$).

P is semifree if $\exists 0 \subset P \subset \mathcal{B} \subset \dots$ exhaustive
filtration by DG-modules such that each quotient is free.
free \implies projective so extensions are semi spl. f. (not spl. f.
for differential!)

$A \subset A^0\text{-DG-mod}$ introduced before is just the DG-category
of finitely generated semi-free DG-modules!

so $\{\text{semi-free modules}\} \xrightarrow{\sim}$ is infinite version of A .

→ $\overset{D}{\rightarrow}$: analog of category of ind-objects (systems $\xrightarrow{\text{in } X_i}$
or contain A^0 -modules .. representable functors)

I abstract category of indices, functor $I \rightarrow A$
→ consider homotopy colimit of this functor DG-object
assoc. to some simplicial object ... colimits of such
I of weak (A^0 -)functors like in $\overset{A}{\rightarrow}$

- inspired by (standard) cell complexes .. ie don't
necessarily attach in correct order.

A DG-algebra. A DG-module P over A^0 is
semi-free, if \exists $0 \subset P_1 \subset P_2 \subset \dots$ ($\cup P_i = P$,
 P_i freely generated by P_{i+1} & homogeneous
generators e.g. $j \in J_i$; so that $d_{i+1} \in P_{i+1}$.
($\& d_{P_i}|_{P_{i+1}} = d_{P_{i+1}}$)

of course P_i/P_{i+1} is free as dg module.

Advantage of above formulation: semi-free modules
sense in nonlinear situations: eg for DG algebras
- same def: exists filtration, of course quotients
won't be algebras . . $P_i = P_{i+1} \langle X_i \rangle$ freely generated
by P_{i+1} & generators X_i , differentials are contained
in subalgebra of old generators.
same for algebras over any operad ..

Eg topological semigroups: attaching cell areas add
(cell in topological sense (maps from sphere),
generate semigroup by old semigroup & this cell with
attaching map. --- construction of algebraic free
generator & topological adding of)

Eg Stasheff operad: to what is it an answer? construct
semifree resolution of the associative operad (in topological operad)

Theorem: \mathbb{E}_1 model structure on A^0 -DG-mod (or in DG or
with W: ~~homotopy equivalence~~, F: surjective
quasi-isomorphisms)

The resulting cofibrant objects (ie morphisms $O \rightarrow P$ are cofibres)
 are retracts of semi-free objects.
 (In a sense : retract means closed summand)

Example 3 A bounded above complex of free R -modules

is semi-free

$\rightarrow P^{n-2} \rightarrow P^{n-1} \rightarrow P_n \rightarrow 0$; $P = \bigoplus P_i$, with stupid filtrations
 --- so bimod above complexes is special case.

Example 4 $\rightarrow \mathbb{Z}/4\mathbb{Z} \xrightarrow{?} \mathbb{Z}/4\mathbb{Z} \xrightarrow{?} \dots$

is not semi-free (can't find basis vector annihilated by the differential!)

Theorem $H_0(\mathcal{A}) \rightarrow D(A)$ is an equivalence.

Proof In bounded above setting: we know 1: anything has bounded above projective resolution Lemma 2: morphism from bounded above projective to acyclic is nullhomotopic.

Lemma 1 $\forall M \in \mathcal{A}^0$ -DG-mod \exists quism $P \rightarrow M$, $P \in \mathcal{A}$
 (in fact can choose surjective - though not necessarily).
 -- a weird category cat!

Lemma 2 $P \in \mathcal{A} \Rightarrow \forall$ acyclic $M \in \mathcal{A}^0$ -DG-mod,
 every morphism $P \rightarrow M$ is homotopic to 0
 -- i.e. P is "K-projective" or homotopically projective.

Exercise M is homotopically projective iff it is homotopy equivalent to a semi-free module.

Proof of Lemma 2 : $P = \bigcup_{i=1}^{\infty} P_i$; P_i/P_{i-1} free . $f: M \rightarrow P$
 i.e. $f \in \text{Hom}(P, M)$ $\deg f = 0$, $\delta f = 0$

Construct boundary: i.e. $G = dH$, $H \in \text{Hom}(P, M)$ $\deg H = -1$
 -- build inductively. $H_i \in \text{Hom}(P_i, M)$ $\deg -1$,
 $dH_i = f_i = f|_{P_i}$.

If given H_{i-1} : choose $\tilde{H}_i|_{P_{i-1}} = H_{i-1}$, \tilde{H}_i deg +
 correct $H_i = \tilde{H}_i + \bar{H}_i$, $\bar{H}_i \in \text{Hom}(P_i/P_{i-1}, M)$ deg +

s.t. $\delta(\tilde{f}_i) = f_i - d\tilde{A}_i \in \text{Ker}(P_i/P_{i-1}, M)$, $d(f_i - d\tilde{A}_i) = 0$

But this DG-module of Hom's is acyclic:

P_i/P_{i-1} is a direct sum of free (shifted) modules, so Ker is a product of shifted copies of $M \Rightarrow$ acyclic \blacksquare

Proof Lemma 1

Sublemma $N \rightarrow M$ morphism of DG-modules

$\Rightarrow \exists$ factorization $N \hookrightarrow N' \rightarrow M$

- s.t.
1. ~~Acyclic~~ $H^*(N') \rightarrow H^*(M)$ bijective
 2. N'/N semi-free

Sublemma \Rightarrow Lemma: construct surjective $f: P_i \rightarrow M$

P_i semi-free & $H^*P_i \rightarrow H^*M$

- Then apply sublemma for $N=P_i \rightarrow P \rightarrow M$

surjective & isom or cohomology:

as M a generator make module with no generators
 e, e' of correct degree, $de = e$, $de' = 0$.

\rightarrow surjectivity on module. similarly for cohomology
 - lift to cocycles.

Proof of sublemma: enough to const-ct $N \hookrightarrow N_i \rightarrow M$, N_i/N semi-free
 with $\text{Ker}(H^*N \rightarrow H^*M) = \text{Ker}(H^*N \rightarrow H^*N_i) \blacksquare$

--- then build $N = UN$: where $N \hookrightarrow N_1 \hookrightarrow N_2 \hookrightarrow \cdots \rightarrow M$

$K := \text{Ker}(H^*N \rightarrow H^*M)$. kick homos generators (H^*A needs)
 ... add cols to fill cohomology

V. Drinfeld - DG Categories III

11/11/02

- 2-category of DG-categories ... partial answer to "What do DG-categories form?" ... reply they form some kind of ∞ -category.

Today: All DG-categories are pretriangulated unless mentioned

Typical example of 2-category: Cat — have objects: categories, morphisms: functors, 2-morphisms: natural transformations

We'll define 2-category DG-cat. Objects: (small) DG-categories
 $A, B \in \text{DG-Cat} \Rightarrow T(A, B)$ category
of quasi-functors. Ob $T(A, B)$ are 1-morphisms of a 2-category,
morphisms in $T(A, B)$ are 2-morphisms of DG-cat.

Composition is functor $T(A, B) \times T(B, C) \rightarrow T(A, C)$

Associativity here is structure not property:

weak associativity: $A \xrightarrow{F} B \xrightarrow{G} C \xrightarrow{\phi} D$
is an isomorphism $((\phi G)F) \xrightarrow{\sim} (\phi(GF))$ functorial
in all three — associativity constraint

→ pentagon axiom for composable 1-morphisms

$$(F_1 F_2) F_3 F_4 = (F_1 (F_2 F_3) F_4)$$

$$(F_1 F_2)(F_3 F_4)$$

$$F_1(F_2 F_3) F_4$$

$$F_1(F_2(F_3 F_4))$$

$\forall A \in \text{DG-Cat} \exists \text{id}_A \in T(A, A)$ s.t.

\exists isomorphism $\text{id}_A \circ F \xrightarrow{\sim} F$ functorial w.r.t F
& similar for $F \circ \text{id}_B \rightarrow F$

Here $T(A, B)$ will in fact be triangulated

- Problem: should really replace $T(A, B)$ by DG-category!
possible over a field, unclear over ring.

$A \in \text{DG-Cat} \mapsto \text{Ho}(A)$ triangulated category

will give a 2-functor $\text{Ho} : \underline{\text{DG-Cat}} \rightarrow \underline{\text{TriagCat}}$

Stupid 2-category : $\underline{\text{DG-Cat}}^{\text{naive}}$: Objects are Obj .
 1-morphisms are DG-functors, 2-morphisms are closed
 morphisms of DG-functors : $F, G : A \rightarrow B$
 a closed morphism is $\gamma : F \rightarrow G$ s.t. $x \in A$ $\gamma_x : F(x) \rightarrow G(x)$
 is closed.

Problems: want to invert quasi-equivalence... also this
 notion of closed morphism not invariant w.r.t. quasi-equivalence.

Ex. An arbitrary DG-category, k ground field construct
 as DG-algebra or DG category with one object
 $T(k, A) (= T(k^{\text{pretr}}, A)) = ?$

Answer = $\text{Ho}(A)$!

Ex. $T(A, k) \xrightarrow[k=\text{inversion}]{\quad}$
 $k \rightarrow k^{\text{pretr}}$
 $k \rightarrow \text{semifree complexes}$

$\left\{ \begin{array}{l} k^{\text{pretr}} : \text{add direct sums and} \\ \text{add zeros for morphisms} \\ \dots \text{ if } k^{\text{pretr}} = \{ \text{finite complexes} \\ \text{of fgpr free} \\ k\text{-modules} \} \end{array} \right.$

Answer: $\overline{T}(A, k) = \text{derived category of } A\text{-modules}$

(recall $A\text{-module is just DG-functor } A \rightarrow \text{complexes} \text{.}$)

How get derived not homotopy category ...

... doesn't change under quasi-equivalence unlike homotopy category

$= \text{Ho}(\xrightarrow{A^\circ})$ left semifree $A\text{-modules}$
 $=: \mathcal{D}(A)$

Ex. $T(A, k^{\text{pretr}}) \subset \mathcal{D}(A)$ full subcategory
 $= \{F : A \rightarrow \text{complexes s.t. } \forall a \in A,$
 $F(a)$ is quasi-isomorphic to a finite complex
 $\text{of free } k\text{-modules of finite rank } G \in \mathcal{D}^{\text{pretr}}\}$

Keller's def $k = \text{field}$

$\xrightarrow{A, B} T(A, B) := \mathcal{D}(A \otimes B^\circ)$ derived category of bimod.
 ... "ind quasi-functors"

$M \in T(A, B)$ gives functor $F_M : \text{Ho}(A) \rightarrow \mathcal{D}(B^\circ)$
 $= \text{Ho}(B)$

$T(A, B) = \{M \in I_*(A, B) : \forall a \in \text{Ho}(A) F_M(a) \text{ is isomorphic to an object of } \text{Ho}(B)\}$

-- can define $DG_{\text{cat}} \subset DG_{\text{cat}}$ where morphisms are ind quasi-functors in $\overrightarrow{I(A, B)} \dots$

Composition: $M \in I(A, B) = D(A \otimes B^\circ)$, $N \in D(B \otimes C^\circ)$
 $\Rightarrow M \overset{L}{\otimes} N \in D(A \otimes C^\circ)$

Note: over a field can lift anything from D to triangulated ... can choose all ass't, tri'nal over a field! M, N semi-free
 then don't need to derive $M \overset{L}{\otimes} N = M \otimes N$ & is semifree.

Not over field: only $\overset{L}{\otimes}$ will be well defined under quasi-equivalence (without see flatness assumptions...)

Any k : everything works under additional assumption of "Purity":

Def A is homotopically flat/ k if all Hom complexes are homotopically flat.

Def A complex C^\bullet is said to be homotopically flat if $C^\bullet \otimes (\text{acyclic})$ is acyclic.

Homotopical Flatness & Projectivity $\text{R}M$ is projective if

$\text{Hom}(M, -)$ is exact,
 complex C^\bullet is homotopically projective if $(fg)_*$ functor $\text{Hom}(P, -)$ preserves acyclicity \Leftrightarrow preserves quasi-isomorphisms

- A semifree complex is homotopically flat (\lim^1 of f_* semifree modules, \otimes commutes with \lim^1 by direct limit)
- Show homotopical projectivity \rightarrow homotopical flatness
- Suppose C is homotopically flat & acyclic \Rightarrow $(C \otimes (\text{any complex}))$ is acyclic ...

sometimes better to have all hom modules semi-free - even better behaved than homotopically flat -

Any DG category has a semi-free resolution, hence in particular a homotopically flat resolution

A DG-category \Rightarrow class of homotopically flat resolutions $\mathcal{F}_i \rightarrow A, i \in I$

- need to compare categories $T(\mathcal{F}_i, B)$

(don't need to resolve with A, B)

$$Ob T(A, B) = \coprod_{\mathcal{F}_i \rightarrow A} T(\mathcal{F}_i, B)$$

- don't have to be strict in defining objects in category since only care about categories up to equivalence

$$\text{or even } \coprod_{\substack{\mathcal{F}_i \rightarrow A \\ \mathcal{B}_j \rightarrow B}} T(\mathcal{F}_i, \mathcal{B}_j).$$

$\forall i, j \in I$ want canonical object $F_{ij} \in T(\mathcal{F}_i, \mathcal{F}_j)$

- "identity on i ", with compatibility

$i, j, k \in I$ want $F_{ij}, F_{jk} \xrightarrow{\sim} F_{ik}$ & inverses.

$$T(\mathcal{F}_i, \mathcal{F}_j) \subset \mathcal{I}(\mathcal{F}_i, \mathcal{F}_j) = D(\widetilde{\mathcal{A}}_i \otimes \mathcal{A}_j^\circ)$$

But $\mathcal{F}_i \xrightarrow{\widetilde{\mathcal{A}}_i} \widetilde{\mathcal{A}}_j \rightarrow$ consider $\mathcal{F}_i \times \mathcal{F}_j^\circ \rightarrow$ complexes
 $\pi_i: \mathcal{F}_i \xrightarrow{\pi_i} A \xleftarrow{\pi_j} \mathcal{F}_j$ $(a \in \mathcal{A}_i, a' \in \mathcal{A}_j) \mapsto \text{Hom}(\pi_i(a), \pi_j(a))$

\Rightarrow canonical F_{ij} .

So this allows us to define DG-cat without imposing flatness, which is often undesirable
- e.g. dg (complexes of abelian groups)

Problem doing all this on DG level:

get not concrete complexes but complexes up to canonical homotopy equivalences \dashrightarrow standard notion
of category, to 2 objects assign many models
of hom complexes, & contractible family of homotopy

equivalences --- even weaker than Ab category ... see
kind of ∞ -category ... so for now
 $T(A, B)$ is only triangulated not DG.

Note: category A is equivalent to
category of pairs $X, Y \in A$ and maps $X \rightarrow Y$,
which is quasi-isomorphism - ie.
homotopy diagonal of A .
Each of the two projectors is a homotopy
equivalence with A .

Kontsevich model

There is a canonical 2-functor $DG_{Cat}^{nab} \xrightarrow{\sim} DG_{Cat}$
(assume only flat versions for simplicity).

$$F: A \otimes B / DG\text{-functor} \Rightarrow \text{bimodule } M_F \in D(A \otimes B^0).$$

analog of bimodule assoc to algebra for $A \leftarrow B$

- namely B is a $(A \otimes B^0)$ -module.
- in categorical version, $M_F: A \times B^0 \rightarrow$ complexes
is $M_F(a, b) = \text{Hom}(b, F(a))$.

$$\text{which actually lies in } T(A, B) \subset I(A, B) = D(A \otimes B^0)$$

and the functor $H_0(A) \rightarrow H_0(B)$ coming from M_F is
just $H_0(F)$

so we have the 2-functor on morphisms

Sasha: any quasi-functor can be viewed as functor
 $A \rightarrow fct B$ -modules which are quasi-representable : }
 $B \nearrow$ quasi to representable are

- ie localization of functors, $A \rightarrow B$...

$$\text{Ext}_{A \otimes B^0}^i(M_F, M_G) = \text{Ext}_{A \otimes B^0}^i(Hom_A, Hom(F, G))$$

Has $Hom(F, G)$ is the DG $A \otimes A^0$ -module
 $(a, a') \mapsto \text{Hom}(F(a'), G(a))$

What is Hom_A ? for algebra A can think
of A as an A -bimodule ...
 Hom_A is the categorical analog: DG- $A \otimes A^{\text{op}}$ -module
 $(a, a') \mapsto \text{Hom}(a; a')$ (i.e. $\text{Hom}(\text{Id}, \text{Id})$)
 gives 2-morphisms for or $\text{DGCat}^{\text{cage}} \rightarrow \text{DGCat}$
 with chain compn.

This gives the full subcategory of
 quasi-functors given by bimodules!
 via $F \xrightarrow{\sim} M_F$ bimodules.
 This subcategory is manifestly self-dual, w/ \mathbb{R}
 Keller equivalence defn.

Ex If A is semi-free then Hom_A of $T(A, B)$ is
 isomorphic to some M_F .

Note $\text{Hom}_A = \text{diagonal bimodule}$

$\text{Ext}^i(\text{Hom}_A, -) = \text{Hoch } i\text{-th cohomology}$
 - gives full description of DG C^k for A
 semi-free by above exercise... in general must
 use semi-free resolutions.
 If A is semi-free, i.e., then morphisms of A are semi-free \Rightarrow
 $\text{Stand}(A)$ has a semi-free resolution
 eg $k = \text{field}$

can compute $T(A, B)$ in terms of this resolution
 $\text{Stand}(A)$:

$\text{Ob}(T(A, B))$ are DG-functors $\text{Stand}(A) \rightarrow \mathbb{R}$.
 aka A_∞-functors $A \rightarrow B$...

$\text{Stand}(A)$ is monad, composition & no adjoint functors,
 maps to the stability ("twisting cochains" of Quillen)
 - gives way to understand A_∞-functors.
 to get DG model this way for $T(A, B)$...

k first Example of comparison of Keller & Kontsevich mod's:

$$T(A, k^{\text{perf}})$$

($k^{\text{perf}} = \text{category with f.d. cohomology}$)

Keller: semifree DG - A-modules M such that
 $\dim_k H^*(M) < \infty$

Kontsevich: finite complexes of fin dim k -vector spaces
with a weak action of A .

Weak action: for any $a \in A$ have an endomorphism $f_a : C \rightarrow C$
(depending linearly on a), $f_{ab} \neq f_a f_b$ necessarily, but

$$f_{ab} - f_a f_b = d(\varphi_{ab}),$$

some combo of φ, f which is cocycle shall be coboundary,
etc...

- come naturally to this by trying to construct f.d. model
of Keller module M . M is homotopy equivalent
to a f.d. complex C , get for $a \in A$ an
endomorphism up to homotopy, will satisfy above relations.

[Gain fin dim by losing strictness of action]

Quotients of DG Categories

Verdier: T triangulated, $Q \subset T$ full triangulated subcategory
get $T \xrightarrow{\quad} T/Q$ quotient triangulated category
with universal property wrt such quotients.

What is DG versn: $B^{\text{full}} \wedge \text{DG catg. rs}$ [Keller]

2-categorical version advantage: works over rings
disadvantage: less precise (weaker form of uniqueness)

Theorem-definition A 2-categorical quotient of A and B is a pair (C, ξ) of a DG category C & a quasifunctor $\xi \in T(A, C)$, s.t. the following equivalent properties hold:

1. $H_0(\xi) : H_0(A) \rightarrow H_0(C)$ identifies $H_0(C)$ with $H_0(A)/H_0(B)$... ie $H_0(A) \xrightarrow{\sim} H_0(C)$ (B killed in H_0) $\xrightarrow{H_0(A)/H_0(B)} S/$ factorization and resulting functor from Verdier quotient is an equivalence).

2. $\forall C' \in \text{DGCat}$ the functor $T(C, C') \rightarrow \text{Ker}(T(A, C') \rightarrow T(B, C'))$ is an equivalence [when $B \xrightarrow{\sim} A \xrightarrow{\sim} C'$ is zero]

(Verdier setting: univariant property $\text{Ker}(T(A, C') \rightarrow T(B, C')) = T(C, C')$ every further killing B factors through C')

Such $(C, \xi) \exists!$ in 2-categorical sense

Not obvious that either implies the other.

Recall $H_0(A) = T(k^{\text{perf}}, A)$ while second property formulated in terms of $T(A, -)$

k field each $T(-)$ is actually a DG-category \Rightarrow more precise version, quivalent in a stronger sense.

Problem is not to prove but to formulate question what k is a ring.

Suppose $T(A, C)$ comes from a DG category $\text{DG}(A, C)$

- above formulation means has collection of objects $\xi_i \in \text{DG}(A, C)$ together with a homotopy class of homotopy equivalences $\xi_i \xrightarrow{\sim} \xi_j$.

- for more precise version need either specific ξ or rather contractible space of morphisms

from the objects ξ_i, ξ_j (ie P resolution of k , $P \rightarrow \text{Hom}(\xi_i, \xi_j)$)

k not a field don't have $\text{DG}(A, C)$ yet...

Def A DG quotient of A and B is a diagram
 $\begin{array}{ccc} A & \xrightarrow{\quad \text{quasi-eq} \quad} & C \\ \downarrow & \text{in particular get object of } & \downarrow \\ T(A, C) & = & T(A, C) \end{array}$

s.t. (A, C) is a 2-categorical quotient.

\leadsto uniqueness (as field) in ∞ -category
of DG-categories (i.e its -model of DG Cat)

... replace T by DG : $DG(C, C') \rightarrow \ker(DG_{SC}) \rightarrow DG_{BC})$
is a quasi-equivalence ~~functor~~, from definition I,
just need to define $DG(CC, C')$

2 constructions. I : new (Drinfeld), works if A
is homotopy flat.

$Ob C = Ob A$, add new morphisms.

$\forall x \in B \subset A$, add $Ex : x \rightarrow x$ of degree -1 ,
 $dEx = id_x$. C is freely generated by A and Ex .

$$A \longrightarrow C$$

If A is homotopy flat this gives desired answer
Rather than inverting quasi-isomorphisms think
of killing objects: i.e killing all morphisms
to & from this object (homotopically)
enough to kill identity morphism?

II. Keller's definition: consider inclusion $\begin{array}{c} A \\ \hookrightarrow \\ B \end{array}$ = semiadditive A -cat
full subcategory: or have
DG functor $M \mapsto M \otimes_A B$ fully faithful or no.

$$A \supset B^\perp = \{M \in A : M|_B \text{ is acyclic}\}$$

$$\hookrightarrow \text{Ext}^1(h_B, M) = 0 \quad \forall b \in B$$

Fact $H_0(B^\perp) \xrightarrow{\sim} H_0(A) / H_0(B)$ h_B representable abelian
is an equivalence. $H_0(A) / H_0(B)$... fully faithful.

So define object $C = \{M \in \mathcal{B}^\perp : \exists a \in A \text{ s.t. } \begin{array}{c} \text{distri.} \\ \text{triangle} \end{array}\}$
 \dashv essential image of $\mathrm{Ho}(C)/\mathrm{Ho}(\mathcal{B}) \xrightarrow{\quad}$

Here don't have DG functor $A \rightarrow C$ but need $\begin{array}{ccc} f & \nearrow & \searrow \\ a & & N \end{array}$
 \dashv try to assign M to a , using a choice,
but with contractible space of choices
(practically equivalent above).

equivalently construct N !
 f of deg 0 ($P = N[-1]$)
so that $\mathrm{core}(f) \in \mathcal{B}^\perp$ i.e. $\mathrm{core}(f)/\mathcal{B}$ is cyclic
i.e. f/\mathcal{B} is a quasi-isomorphism

Restriction to \mathcal{B} : think of a via representable functor $\mathrm{h}_{\mathcal{B}} \circ {}^0\mathrm{-DG-mod}$
take $\mathrm{h}_{\mathcal{B}}$ and take a semi-free resolution $P \rightarrow \mathrm{h}_{\mathcal{B}}$
— this is our choice, but unique
in strongest homotopical sense (contractible space of choices).

Not self-dual definition... so would like universal property
to identify all.

How might you come to this def? why
semi-free Ob -modules? if believe Ob quotient
exists: $A \rightarrow C$ with property 1

Yoneda: $C \hookrightarrow C^0\text{-DG-mod} \xrightarrow{\quad} A^0\text{-DG-mod}$
 $\Rightarrow \mathrm{h}_C(C) \rightarrow \mathrm{h}_A(A^0\text{-DG-mod}) \xrightarrow{\text{restrict.}} \mathrm{D}(A^0) = \mathrm{h}_A(A)$
- fully faithful, so natural to look for C in A .

In fact Δ (infinite version) is implicit in Verdier:

$\mathrm{Ob} \mathcal{T}/Q = \mathrm{Ob} \mathcal{T}$, Morphisms: invert Q -quasi-isomorphisms
- i.e. morphisms with core in Q .

So morphisms are zig-zags $\square \leftarrow \square \rightarrow \square$, $\square \leftarrow Q \rightarrow \square$
but subject to take $\square' \leftarrow \square \rightarrow \square$ or $\square \leftarrow \square' \rightarrow \square$.

Verder: $X, Y \in \text{Ob } T$ $\text{Ext}_{T/\mathcal{P}}^i(X, Y) =$
 $= \varprojlim_{\mathcal{C}_Y} \text{Ext}_T^i(X, Y)$ over all "injective resolutions" $Y \rightarrowtail Y'$,
 $\mathcal{C}_Y = \text{category with objects } Q\text{-quivers } Y \rightarrow Y'$

ie maps $X \rightarrow Y'$ give map $X \rightarrow Y$ in said category
 \mathcal{C}_Y is a filtering category.
 \Rightarrow --- need for group structure on the limit of Ext's
 (filtering: commute with finite projective limits (eg products))

- claim this is same infinity as \varprojlim^A .
- can write via homotopy colimit, but Keller's construction takes care of this via semi-fineness.

Exercise $a, a' \in A$, $P \xrightarrow{\text{qis}} \mathbb{A}_a / \mathbb{A}_B$ semi-free resolution
 $\xrightarrow{P \in \mathcal{B}}$. Show directly that

$$\text{Ext}_{H_0(A)/H_0(B)}^i(a', a) = \varprojlim_M H^i \text{Hom}(a', M) \quad \text{Verder}$$

$\xrightarrow{M \in \{M, S \rightarrow M\}}$
 \parallel with coe in $H_0(\mathcal{B})\}$

Keller answer $H^i(\text{Hom}_A(a', \text{cone}(P \otimes_B \xrightarrow{\sim} a)))$ up to homotopy

- write $P = \bigvee P_i$ union of fg modules (filtering family)
- ie cohomology of injective limit over filtering vs direct limit of cohomology -- slightly different filtering categories...
- Get from a' to a by cofinality of \mathcal{A} so quasi-equivalence it is an equivalence.